

# Basic Data Structures in Python

Stack, Queue

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Class

# Class Definitions i

A class serves as the primary means for abstraction in object-oriented programming. In Python, every piece of data is represented as an instance of some class. A class

- provides a set of behaviors in the form of **member functions** (also known as methods), with implementations that are common to all instances of that class.
- serves as a blueprint for its instances, effectively determining the way that state information for each instance is represented in the form of attributes (also known as **fields, instance variables, or data members**).

# Class

Example: CreditCard Class

## Example: CreditCard Class i

```
1 class CreditCard:
2     """A consumer credit card."""
3     def __init__(self, customer, bank, acct, limit):
4         """Create a new credit card instance.
5
6         The initial balance is zero.
7
8         customer    the name of the customer (e.g., John Bowman )
9         bank        the name of the bank (e.g., California Savings )
10        acct        the account identifier (e.g., 5391 0375 9387 5309 )
11        limit       credit limit (measured in dollars)
12    """
13    self.customer = customer
14    self.bank = bank
15    self.account = acct
16    self.limit = limit
17    self.balance = 0
```

## Example: CreditCard Class ii

```
18
19 def get_customer(self):
20     """Return name of the customer."""
21     return self.customer
22
23 def get_bank(self):
24     """Return the bank s name."""
25     return self.bank
26
27 def get_account(self):
28     """Return the card identifying number (typically stored as a string)"""
29     return self.account
30
31 def get_limit(self):
32     """Return current credit limit."""
33     return self.limit
34
35 def get_balance(self):
```

## Example: CreditCard Class iii

```
36     """Return current balance."""
37     return self.balance
38
39     def charge(self, price):
40         """Charge given price to the card, assuming sufficient credit limit.
41
42         Return True if charge was processed; False if charge was denied."""
43
44         # if charge would exceed limit,
45         if price + self.balance > self.limit:
46             return False # cannot accept charge
47         else:
48             self.balance += price
49             return True
50
51     def make_payment(self, amount):
52         """Process customer payment that reduces balance."""
53         self.balance -= amount
```



## Example: CreditCard Class iv

```
54
55
56 if name == '__main__':
57     wallet = [ ]
58     wallet.append(CreditCard( John Bowman , California Savings ,
59                               5391 0375 9387 5309 , 2500) )
60     wallet.append(CreditCard( John Bowman , California Federal ,
61                               3485 0399 3395 1954 , 3500) )
62     wallet.append(CreditCard( John Bowman , California Finance ,
63                               5391 0375 9387 5309 , 5000) )
64
65     for val in range(1, 17):
66         wallet[0].charge(val)
67         wallet[1].charge(2 * val)
68         wallet[2].charge(3 * val)
69
70     for c in range(3):
71         print( Customer = , wallet[c].get_customer())
```

## Example: CreditCard Class v

```
72     print( Bank = , wallet[c].get bank())
73     print( Account = , wallet[c].get account())
74     print( Limit = , wallet[c].get limit())
75     print( Balance = , wallet[c].get balance())
76     while wallet[c].get balance( ) > 100:
77         wallet[c].make payment(100)
78         print( New balance = , wallet[c].get balance())
79     print( )
```

Stack

## Definition

Stack A stack is a collection of objects that are inserted and removed according to the last-in, first-out (LIFO) principle.

A user may insert objects into a stack at any time, but may only access or remove the most recently inserted object that remains (at the so-called “top” of the stack).

### Example

Internet Web browsers store the addresses of recently visited sites in a stack. Each time a user visits a new site, that site's address is "pushed" onto the stack of addresses. The browser then allows the user to "pop" back to previously visited sites using the "back" button.

### Example

Text editors usually provide an "undo" mechanism that cancels recent editing operations and reverts to former states of a document. This undo operation can be accomplished by keeping text changes in a stack.

# Stack

The Stack Abstract Data Type

# The Stack Abstract Data Type i

Stacks are the simplest of all data structures, yet they are also among the most important.

Formally, a stack is an abstract data type (ADT) such that an instance  $S$  supports the following two methods:

- $S.\text{push}(e)$ : Add element  $e$  to the top of stack  $S$ .
- $S.\text{pop}()$ : Remove and return the top element from the stack  $S$ ; an error occurs if the stack is empty.

## The Stack Abstract Data Type ii

Additionally, define the following accessor methods for convenience:

- `S.top()`: Return a reference to the top element of stack `S`, without removing it; an error occurs if the stack is empty.
- `S.is_empty()`: Return `True` if stack `S` does not contain any elements.
- `len(S)`: Return the number of elements in stack `S`; in Python, we implement this with the special method `__len__`.



# The Stack Abstract Data Type iii

**Example 6.3:** *The following table shows a series of stack operations and their effects on an initially empty stack S of integers.*

Operation	Return Value	Stack Contents
S.push(5)	–	[5]
S.push(3)	–	[5, 3]
len(S)	2	[5, 3]
S.pop()	3	[5]
S.is_empty()	False	[5]
S.pop()	5	[]
S.is_empty()	True	[]
S.pop()	“error”	[]
S.push(7)	–	[7]
S.push(9)	–	[7, 9]
S.top()	9	[7, 9]
S.push(4)	–	[7, 9, 4]
len(S)	3	[7, 9, 4]
S.pop()	4	[7, 9]
S.push(6)	–	[7, 9, 6]
S.push(8)	–	[7, 9, 6, 8]
S.pop()	8	[7, 9, 6]

# Stack

Simple Array-Based Stack Implementation

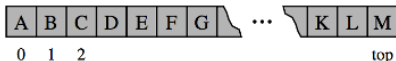
# Simple Array-Based Stack Implementation i

We can implement a stack quite easily by storing its elements in a Python list.

The list class already supports

- adding an element to the end with the **append** method,
- removing the last element with the **pop** method,

so it is natural to align the top of the stack at the end of the list, as shown in



**Figure 6.2:** Implementing a stack with a Python list, storing the top element in the rightmost cell.

## Simple Array-Based Stack Implementation ii

Although a programmer could directly use the list class in place of a formal stack class,

- lists also include behaviors (e.g., adding or removing elements from arbitrary positions) that would break the abstraction that the stack ADT represents.
- the terminology used by the list class does not precisely align with traditional nomenclature for a stack ADT, in particular the distinction between **append** and **push**.

# The Adapter Pattern i

## Definition

The adapter design pattern applies to any context where we effectively want to modify an existing class so that its methods match those of a related, but different, class or interface.

One general way to apply the adapter pattern is to define a new class in such a way that it contains an [instance of the existing class as a hidden field](#), and then to implement each method of the new class using methods of this hidden instance variable.

<i>Stack Method</i>	<i>Realization with Python list</i>
S.push(e)	L.append(e)
S.pop()	L.pop()
S.top()	L[-1]
S.is_empty()	len(L) == 0
len(S)	len(L)

**Table 6.1:** Realization of a stack S as an adaptation of a Python list L.

# Implementing a Stack Using a Python List i

We use the adapter design pattern to define an `ArrayStack` class that uses an underlying Python list for storage.

One question that remains is what our code should do if a user calls `pop` or `top` when the stack is empty. Our ADT suggests that an error occurs, but we must decide what type of error.

```
1 class Empty( Exception ):
2     pass
```

# Implementing a Stack Using a Python List ii

```
1  class ArrayStack:
2      def __init__(self):
3          self._data = []
4
5      def __len__(self):
6          return len(self._data)
7
8      def is_empty(self):
9          return len(self._data) == 0
10
11     def push(self, e):
12         self._data.append(e)
13
14     def pop(self):
15         if self.is_empty():
16             print('Stack is empty!')
17         return self._data.pop()
18
```

# Implementing a Stack Using a Python List iii

```
19     def top(self):
20         if self.is_empty():
21             print('Stack is empty')
22             return self._data[-1]
23
24     if __name__ == "__main__":
25         S = ArrayStack()
26         S.push(5)
27         S.push(3)
28         print(S._data)
29         print(S.pop())
30         print(S.is_empty())
31         print(S.pop())
32         print(S.is_empty())
33         S.push(7)
34         S.push(9)
35         S.push(4)
36         print(S.pop())
```



# Implementing a Stack Using a Python List iv

```
37 S.push(6)
38 S.push(8)
39 print(S._data)
```

# Implementing a Stack Using a Python List v

The result of running the code

```
1 [5, 3] 3 False
2 5 True
3 4
4 [7, 9, 6, 8]
```

# Implementing a Stack Using a Python List vi

Operation	Running Time
<code>S.push(e)</code>	$O(1)^*$
<code>S.pop()</code>	$O(1)^*$
<code>S.top()</code>	$O(1)$
<code>S.is_empty()</code>	$O(1)$
<code>len(S)</code>	$O(1)$

\*amortized

**Table 6.2:** Performance of our array-based stack implementation. The bounds for push and pop are amortized due to similar bounds for the list class. The space usage is  $O(n)$ , where  $n$  is the current number of elements in the stack.

The  $O(1)$  time for push and pop are amortized bounds.

- A typical call to either of these methods uses constant time;
- But there is occasionally an  $O(n)$ -time worst case, where  $n$  is the current number of elements in the stack, when an operation causes the list to resize its internal array.

# Stack Application

# Reversing Data Using a Stack i

As a consequence of the LIFO protocol, a stack can be used as a general tool to reverse a data sequence.

## Example

If the values 1, 2, and 3 are pushed onto a stack in that order, they will be popped from the stack in the order 3, 2, and then 1.

## Example

We might wish to print lines of a file in reverse order in order to display a data set in decreasing order rather than increasing order.

# Reversing Data Using a Stack ii

## code/reverse.py

```
1 from array_stack import ArrayStack
2
3 def reverse_file(filename):
4     S = ArrayStack()
5
6     original = open(filename)
7     for line in original:
8         S.push(line.rstrip('\n'))
9     original.close
10
11     output = open(filename, 'w')
12     while not S.is_empty():
13         output.write(S.pop() + '\n')
14     output.close()
15
16 if __name__ == "__main__":
17     reverse_file('text.txt')
```

## before

```
this is sixth line
this is fifth line
this is fourth line
this is third line
this is second line
this is first line
```

## after

```
this is first line
this is second line
this is third line
this is fourth line
this is fifth line
this is sixth line
```

# Matching Parentheses i

Consider arithmetic expressions that may contain various pairs of grouping symbols

```
1  from array_stack import ArrayStack
2  def is_matched_html(raw):
3      S = ArrayStack()
4      j = 0
5      while j != -1:
6          k = raw.find('>', j+1)
7          if k == -1:
8              return False
9          tag = raw[j+1:k]
10         if not tag.startswith('/'):
11             S.push(tag)
12         else:
13             if S.is_empty():
14                 return False
```

# Matching Parentheses ii

```
15     if tag[1:] != S.pop():
16         return False
17     j = raw.find('<', k+1)
18     return S.is_empty()
19
20 if __name__ == "__main__":
21     raw = "<a> <b> </b> </a>"
22     if(is_matched_html(raw)):
23         print("Matched")
24     else:
25         print("Not matching")
```



# Matching Tags in HTML i

Another application of matching delimiters is in the validation of markup languages such as HTML or XML.

HTML is the standard format for hyperlinked documents on the Internet and XML is an extensible markup language used for a variety of structured data sets.

## Matching Tags in HTML ii

In an HTML document, portions of text are delimited by HTML tags. A simple opening and corresponding closing HTML tag has the form

```
1 <name> ... </name>
```

Other commonly used HTML tags that are used in this example include:

- body: document body
- h1: section header
- center: center justify
- p: paragraph
- ol: numbered (ordered) list
- li: list item

## Matching Tags in HTML iii

```
1 from array_stack import ArrayStack
2 def is_matched_html(raw):
3     S = ArrayStack()
4     j = 0
5     while j != -1:
6         k = raw.find('>', j+1)
7         if k == -1:
8             return False
9         tag = raw[j+1:k]
10        if not tag.startswith('/'):
11            S.push(tag)
12        else:
13            if S.is_empty():
14                return False
15            if tag[1:] != S.pop():
16                return False
17        j = raw.find('<', k+1)
18    return S.is_empty()
```

# Matching Tags in HTML iv

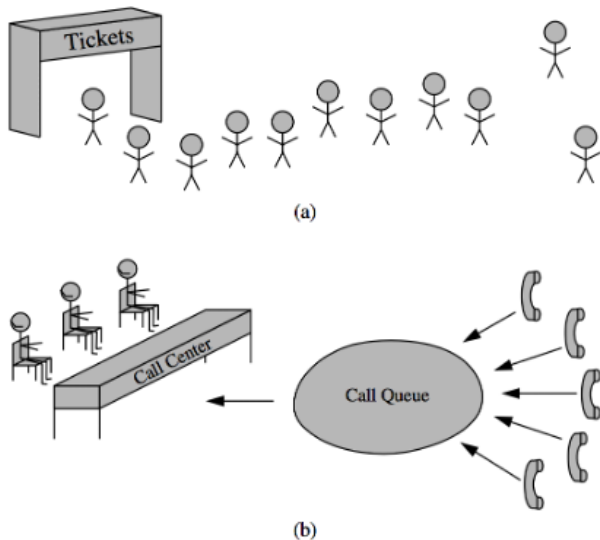
```
19
20 if __name__ == "__main__":
21     raw = "<a> <b> </b> </a>"
22     if(is_matched_html(raw)):
23         print("Matched")
24     else:
25         print("Not matching")
```

Queue

## **Definition**

A queue is a collection of objects that are inserted and removed according to the first-in, first-out (FIFO) principle.

Elements can be inserted at any time, but only the element that has been in the queue the longest can be next removed.



**Figure 6.4:** Real-world examples of a first-in, first-out queue. (a) People waiting in line to purchase tickets; (b) phone calls being routed to a customer service center.

# Queue

The Queue Abstract Data Type



# The Queue Abstract Data Type i

Formally, the queue abstract data type defines a collection that keeps objects in a sequence, where

- element access and deletion are restricted to the first element in the queue;
- and element insertion is restricted to the back of the sequence

# The Queue Abstract Data Type ii

The queue abstract data type (ADT) supports the following two fundamental methods for a queue  $Q$ :

- $Q.enqueue(e)$ : Add element  $e$  to the back of queue  $Q$ .
- $Q.dequeue()$ : Remove and return the first element from queue  $Q$ ; an error occurs if the queue is empty.

The queue ADT also includes the following supporting methods

- $Q.first()$ : Return a reference to the element at the front of queue  $Q$ , without removing it; an error occurs if the queue is empty.
- $Q.is\_empty()$ : Return True if queue  $Q$  does not contain any elements.
- $len(Q)$ : Return the number of elements in queue  $Q$ ; in Python, we implement this with the special method `__len__`.

## The Queue Abstract Data Type iii

By convention, we assume that a newly created queue is empty, and that there is no a priori bound on the capacity of the queue. Elements added to the queue can have arbitrary type.

# The Queue Abstract Data Type iv

**Example 6.4:** *The following table shows a series of queue operations and their effects on an initially empty queue  $Q$  of integers.*

Operation	Return Value	first $\leftarrow$ $Q$ $\leftarrow$ last
Q.enqueue(5)	–	[5]
Q.enqueue(3)	–	[5, 3]
len(Q)	2	[5, 3]
Q.dequeue()	5	[3]
Q.is_empty()	False	[3]
Q.dequeue()	3	[ ]
Q.is_empty()	True	[ ]
Q.dequeue()	“error”	[ ]
Q.enqueue(7)	–	[7]
Q.enqueue(9)	–	[7, 9]
Q.first()	7	[7, 9]
Q.enqueue(4)	–	[7, 9, 4]
len(Q)	3	[7, 9, 4]
Q.dequeue()	7	[9, 4]

# Queue

## Array-Based Queue Implementation

## Array-Based Queue Implementation i

For the stack ADT, we created a very simple adapter class that used a Python list as the underlying storage. It may be very tempting to use a similar approach for supporting the queue ADT.

- We could enqueue element `e` by calling `append(e)` to add it to the end of the list.
- We could use the syntax `pop(0)`, as opposed to `pop()`, to intentionally remove the first element from the list when dequeuing.

As easy as this would be to implement, it is **tragically inefficient**.

## Array-Based Queue Implementation ii

As discussed before, when `pop` is called on a list with a non-default index, a loop is executed to shift all elements beyond the specified index to the left, so as to fill the hole in the sequence caused by the `pop`. Therefore, a call to `pop(0)` always causes the worst-case behavior of  $O(n)$  time.

## Array-Based Queue Implementation iii

We can improve on the above strategy by avoiding the call to `pop(0)` entirely.

We can

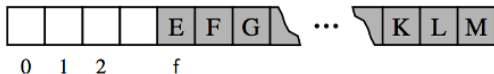
- replace the dequeued entry in the array with a reference to `None`, and
- maintain an explicit variable `f` to store the index of the element that is currently at the front of the queue.

Such an algorithm for dequeue would run in  $O(1)$  time.



## Array-Based Queue Implementation iv

After several dequeue operations, this approach might lead to



**Figure 6.5:** Allowing the front of the queue to drift away from index 0.

## Array-Based Queue Implementation v

Unfortunately, there remains a drawback to the revised approach.

We can build a queue that has relatively few elements, yet which are stored in an arbitrarily large list. This occurs, for example, if we repeatedly enqueue a new element and then dequeue another (allowing the front to drift rightward). Over time, the size of the underlying list would grow to  $O(m)$  where  $m$  is the total number of enqueue operations since the creation of the queue, rather than the current number of elements in the queue.

## Using an Array Circularly i

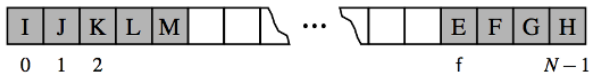
In developing a more robust queue implementation, we allow

- the front of the queue to drift rightward,
- the contents of the queue to “wrap around” the end of an underlying array.

## Using an Array Circularly ii

We assume that our **underlying array has fixed length  $N$**  that is greater than the actual number of elements in the queue.

New elements are enqueued toward the “end” of the current queue, progressing from the front to index  $N-1$  and continuing at index 0, then 1.



**Figure 6.6:** Modeling a queue with a circular array that wraps around the end.

Implementing this circular view is not difficult.

- When we dequeue an element and want to “advance” the front index, we use the arithmetic  $f = (f + 1) \% N$ .

## Using an Array Circularly iii

Internally, the queue class maintains the following three instance variables:

- `_data`: is a reference to a list instance with a fixed capacity.
- `_size`: is an integer representing the current number of elements stored in the queue (as opposed to the length of the data list).
- `_front`: is an integer that represents the index within data of the first element of the queue (assuming the queue is not empty).

# A Python Queue Implementation i

```
1 class ArrayQueue:
2     DEFAULT_CAPACITY = 10
3     def __init__(self):
4         self._data = [None] * ArrayQueue.DEFAULT_CAPACITY
5         self._size = 0
6         self._front = 0
7
8     def __len__(self):
9         return self._size
10
11    def is_empty(self):
12        return self._size == 0
13
14    def first(self):
15        if self.is_empty():
16            print('Queue is empty')
17        return self._data[self._front]
```

# A Python Queue Implementation ii

```
18
19 def dequeue(self):
20     if self.is_empty():
21         print('Queue is empty')
22     result = self._data[self._front]
23     self._front = (self._front + 1) % len( self._data)
24     self._size -= 1
25     return result
26
27 def enqueue(self, e):
28     if self.size == len(self.data):
29         self.resize(2*len(self.data))
30     avail = (self.front + self.size) % len(self.data)
31     self.data[avail] = e
32     self.size += 1
33
34 def _resize(self, cap):
35     old = self.data
```

## A Python Queue Implementation iii

```
36 self._data = [None] * cap
37 walk = self._front
38 for k in range(self._size)
39     self._data[k] = old[walk]
40     walk = (1 + walk) % len(old)
41 self._front = 0
```

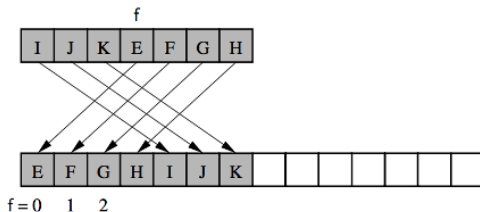


## Resizing the Queue i

When enqueue is called at a time when the size of the queue equals the size of the underlying list, we rely on a standard technique of doubling the storage capacity of the underlying list.

After creating a temporary reference to the old list of values, we allocate a new list that is twice the size and copy references from the old list to the new list. While transferring the contents, we intentionally realign the front of the queue with index 0 in the new array, as shown in

## Resizing the Queue ii



**Figure 6.7:** Resizing the queue, while realigning the front element with index 0.

With the exception of the resize utility, all of the methods rely on a constant number of statements involving arithmetic operations, comparisons, and assignments. Therefore, each method runs in worst-case  $O(1)$  time, except for **enqueue** and **dequeue**, which have amortized bounds of  $O(1)$  time

Operation	Running Time
Q.enqueue(e)	$O(1)^*$
Q.dequeue()	$O(1)^*$
Q.first()	$O(1)$
Q.is_empty()	$O(1)$
len(Q)	$O(1)$

\*amortized

**Table 6.3:** Performance of an array-based implementation of a queue. The bounds for enqueue and dequeue are amortized due to the resizing of the array. The space usage is  $O(n)$ , where  $n$  is the current number of elements in the queue.

# Double-Ended Queues

# Double-Ended Queues i

## Definition

A **dequeue** (i.e., **double-ended queue**) is a queue-like data structure that supports insertion and deletion at both the front and the back of the queue.

**Deque** is usually pronounced “**deck**” to avoid confusion with the **dequeue** method of the regular queue ADT, which is pronounced like the abbreviation “**D.Q.**”.

The deque abstract data type is more general than both the stack and the queue ADTs.

# Double-Ended Queues ii

## Example

A restaurant using a queue to maintain a waitlist.

- Occasionally, the first person might be removed from the queue only to find that a table was not available; typically, the restaurant will re-insert the person at the first position in the queue.
- It may also be that a customer at the end of the queue may grow impatient and leave the restaurant.

# Double-Ended Queues

The Deque Abstract Data Type

# The Deque Abstract Data Type i

To provide a symmetrical abstraction, the deque ADT is defined so that deque D supports the following methods:

- `D.add_first(e)`: Add element e to the front of deque D.
- `D.add_last(e)`: Add element e to the back of deque D.
- `D.delete_first()`: Remove and return the first element from deque D; an error occurs if the deque is empty.
- `D.delete_last()`: Remove and return the last element from deque D; an error occurs if the deque is empty.



# The Deque Abstract Data Type ii

Additionally, the deque ADT will include the following accessors:

- `D.first()`: Return (but do not remove) the first element of deque D; an error occurs if the deque is empty.
- `D.last()`: Return (but do not remove) the last element of deque D; an error occurs if the deque is empty.
- `D.is_empty()`: Return True if deque D does not contain any elements.
- `len(D)`: Return the number of elements in deque D; in Python, we implement this with the special method `__len__`

## The Deque Abstract Data Type iii

**Example 6.5:** *The following table shows a series of operations and their effects on an initially empty deque D of integers.*

<b>Operation</b>	<b>Return Value</b>	<b>Deque</b>
D.add_last(5)	–	[5]
D.add_first(3)	–	[3, 5]
D.add_first(7)	–	[7, 3, 5]
D.first()	7	[7, 3, 5]
D.delete_last()	5	[7, 3]
len(D)	2	[7, 3]
D.delete_last()	3	[7]
D.delete_last()	7	[]
D.add_first(6)	–	[6]
D.last()	6	[6]
D.add_first(8)	–	[8, 6]
D.is_empty()	False	[8, 6]
D.last()	6	[8, 6]

# Double-Ended Queues

Implementing a Deque with a Circular Array

# Implementing a Deque with a Circular Array i

We can implement the deque ADT in much the same way as the ArrayQueue class.

- We recommend maintaining the same three instance variables: **`_data`, `_size`, and `_front`**.
- Whenever we need to know the index of the **back of the deque, or the first available slot beyond the back of the deque**, we use modular arithmetic for the computation.

- in **`last()`** method, uses the index

```
back = (self._front + self._size - 1) % len(self._data)
```

- in **`add_first()`** method, circularly decrement the index

```
self._front = (self._front - 1) % len(self._data)
```

# Double-Ended Queues

Dequeues in the Python Collections Module

## Dequeues in the Python Collections Module i

An implementation of a deque class is available in Python's standard collections module. A summary of the most commonly used behaviors of the `collections.deque` class is given in

## Dequeues in the Python Collections Module ii

<b>Our Deque ADT</b>	<b>collections.deque</b>	Description
len(D)	len(D)	number of elements
D.add_first()	D.appendleft()	add to beginning
D.add_last()	D.append()	add to end
D.delete_first()	D.popleft()	remove from beginning
D.delete_last()	D.pop()	remove from end
D.first()	D[0]	access first element
D.last()	D[-1]	access last element
	D[j]	access arbitrary entry by index
	D[j] = val	modify arbitrary entry by index
	D.clear()	clear all contents
	D.rotate(k)	circularly shift rightward k steps
	D.remove(e)	remove first matching element
	D.count(e)	count number of matches for e

**Table 6.4:** Comparison of our deque ADT and the collections.deque class.

The image features a series of concentric circles in shades of red and orange, creating a tunnel-like effect that leads to a dark blue circular center. The text "That's all Folks!" is written in a white, cursive font across the center.

*That's all Folks!*